



**SRI AKILANDESWARI WOMEN'S COLLEGE, WANDIWASH**

**DESIGN AND ANALYSIS OF ALGORITHM**

**Class : III UG Computer Science**

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## WHAT IS AN ALGORITHM?

- An algorithm is “a finite set of precise instructions for performing a computation or for solving a problem”
  - A program is one type of algorithm
    - All programs are algorithms
    - Not all algorithms are programs!
  - Directions to somebody’s house is an algorithm
  - A recipe for cooking a cake is an algorithm
  - The steps to compute the cosine of  $90^\circ$  is an algorithm

# Some algorithms are harder than others

- Some algorithms are easy
  - Finding the largest (or smallest) value in a list
  - Finding a specific value in a list
- Some algorithms are a bit harder
  - Sorting a list
- Some algorithms are very hard
  - Finding the shortest path between Miami and Seattle
- Some algorithms are essentially impossible
  - Factoring large composite numbers

## ALGORITHM 1: MAXIMUM ELEMENT

- Given a list, how do we find the maximum element in the list?
- To express the algorithm, we'll use pseudocode
  - Pseudocode is kinda like a programming language, but not really

## ALGORITHM 1: MAXIMUM ELEMENT

- Algorithm for finding the maximum element in a list:

**procedure** max ( $a_1, a_2, \dots, a_n$ : integers)

$max := a_1$

**for**  $i := 2$  **to**  $n$

**if**  $max < a_i$  **then**  $max := a_i$

{  $max$  is the largest element }

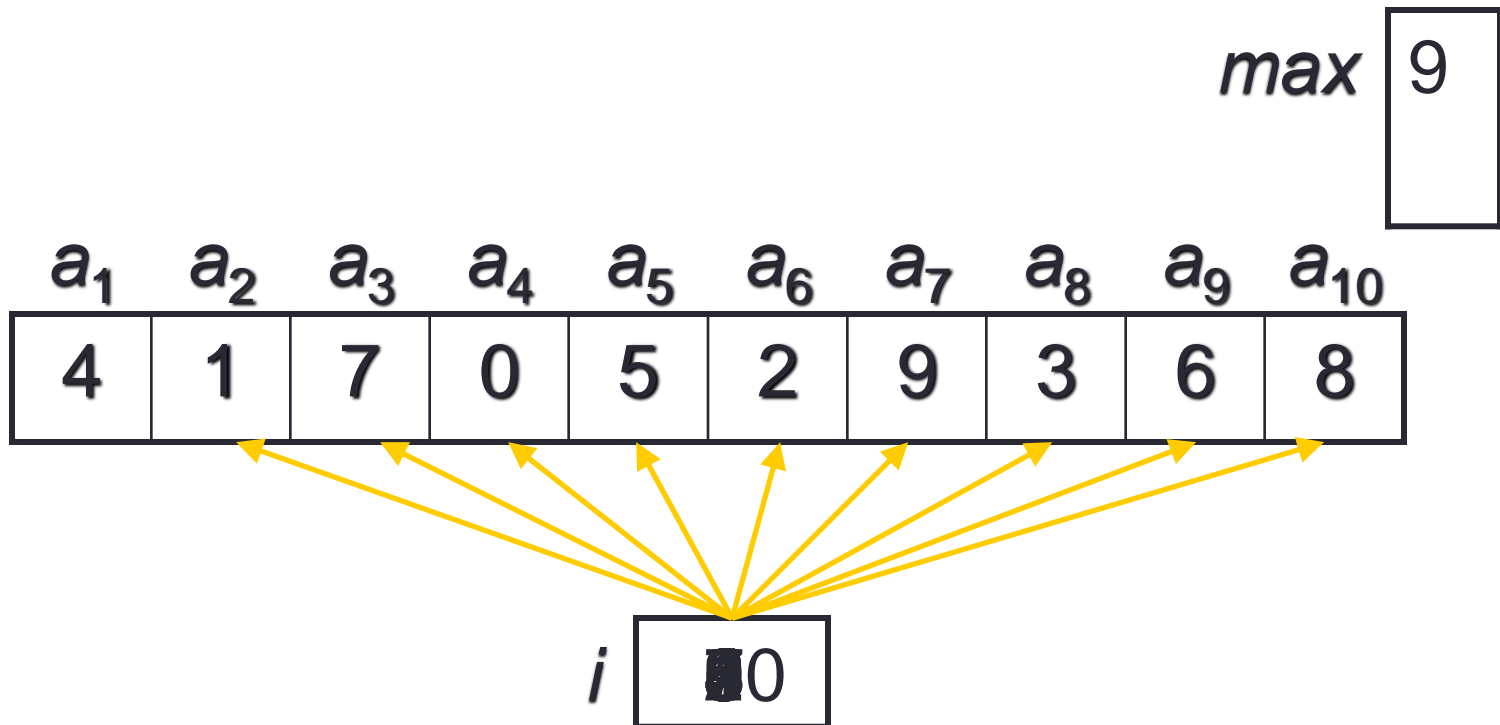
# ALGORITHM 1: MAXIMUM ELEMENT

**procedure** max ( $a_1, a_2, \dots, a_n$ : integers)

$max := a_1$

**for**  $i := 2$  to  $n$

**if**  $max < a_i$  **then**  $max := a_i$



# PROPERTIES OF ALGORITHMS

- Algorithms generally share a set of properties:
  - Input: what the algorithm takes in as input
  - Output: what the algorithm produces as output
  - Definiteness: the steps are defined precisely
  - Correctness: should produce the correct output
  - Finiteness: the steps required should be finite
  - Effectiveness: each step must be able to be performed in a finite amount of time
  - Generality: the algorithm *should* be applicable to all problems of a similar form

## ALGORITHM 2: LINEAR SEARCH

- Given a list, find a specific element in the list
  - List does NOT have to be sorted!

**procedure** linear\_search ( $x$ : integer;  $a_1, a_2, \dots, a_n$ : integers)

$i := 1$

**while** (  $i \leq n$  and  $x \neq a_i$  )

$i := i + 1$

**if**  $i \leq n$  **then**  $location := i$

**else**  $location := 0$

{  $location$  is the subscript of the term that equals  $x$ , or it is 0 if  $x$  is not found }



## ALGORITHM 2: LINEAR SEARCH, TAKE 2

**procedure** linear\_search ( $x$ : integer;  $a_1, a_2, \dots, a_n$ : integers)

$i := 1$

**while** ( $i \leq n$  and  $x \neq a_i$ )

$i := i + 1$

**if**  $i \leq n$  **then**  $location := i$

**else**  $location := 0$

$x$ 

11
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$location$ 

0
---

$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	$a_8$	$a_9$	$a_{10}$
4	1	7	0	5	2	9	3	6	8

$i$ 

10
----

## ALGORITHM 3: BINARY SEARCH

```
procedure binary_search ( $x$ : integer;  $a_1, a_2, \dots, a_n$ : increasing
  integers)
 $i := 1$            {  $i$  is left endpoint of search interval }
 $j := n$           {  $j$  is right endpoint of search interval }
while  $i < j$ 
begin
   $m := \lfloor (i+j)/2 \rfloor$       {  $m$  is the point in the middle }
  if  $x > a_m$  then  $i := m+1$ 
  else  $j := m$ 
end
if  $x = a_i$  then  $location := i$ 
else  $location := 0$ 
{  $location$  is the subscript of the term that equals  $x$ , or it is 0 if  $x$ 
  is not found }
```

## Algorithm 3: Binary search, take 1

**procedure** `binary_search` (`x`: integer; `a`<sub>1</sub>, `a`<sub>2</sub>, ..., `a`<sub>*n*</sub>: increasing integers)

`i` := 1

`j` := `n`

**while** `i` < `j`

**begin**

`m` :=  $\lfloor (i+j)/2 \rfloor$

**if** `x` > `a`<sub>*m*</sub> **then** `i` := `m` + 1

**else** `j` := `m`

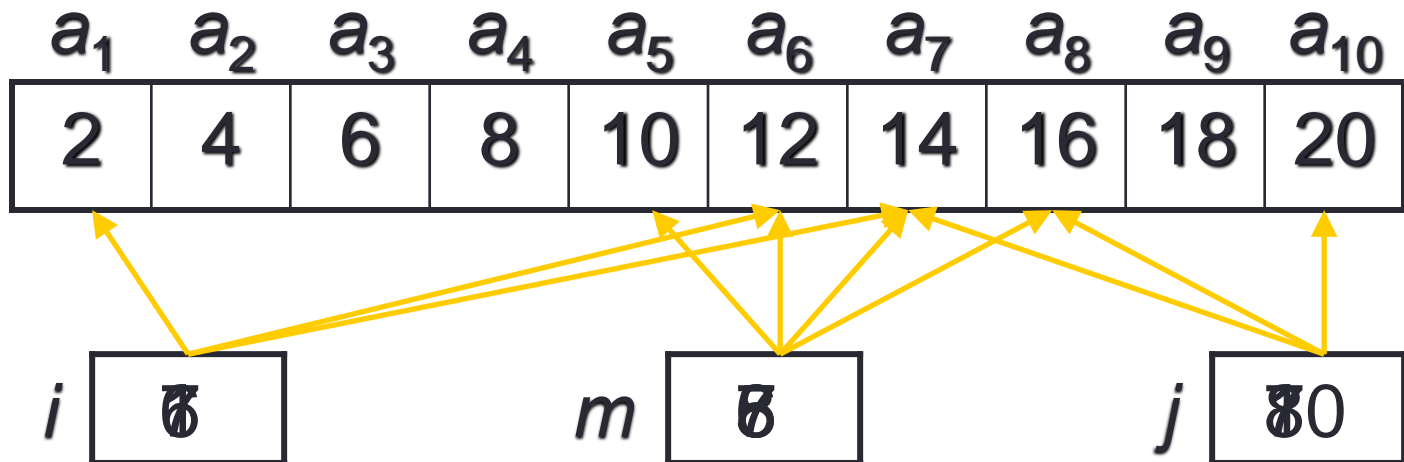
**end**

**if** `x` = `a`<sub>*i*</sub> **then** `location` := `i`

**else** `location` := 0

`x` 14

`location` 7



## ALGORITHM 3: BINARY SEARCH, TAKE 2

**procedure** `binary_search` (`x`: integer; `a`<sub>1</sub>, `a`<sub>2</sub>, ..., `a`<sub>*n*</sub>: increasing integers)

`i` := 1

`j` := `n`

**while** `i` < `j`

**begin**

`m` :=  $\lfloor (i+j)/2 \rfloor$

**if** `x` > `a`<sub>*m*</sub> **then** `i` := `m` + 1

**else** `j` := `m`

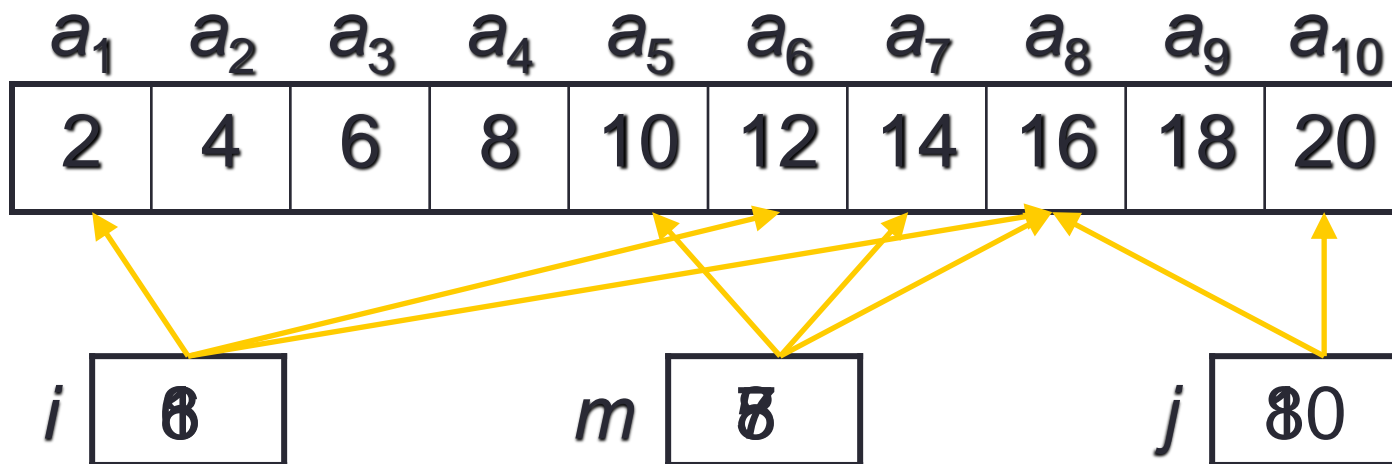
**end**

**if** `x` = `a`<sub>*l*</sub> **then** `location` := *l*

**else** `location` := 0

`x` 15

`location` 0



# BINARY SEARCH RUNNING TIME

- How long does this take (worst case)?
- If the list has 8 elements
  - It takes 3 steps
- If the list has 16 elements
  - It takes 4 steps
- If the list has 64 elements
  - It takes 6 steps
- If the list has  $n$  elements
  - It takes  $\log_2 n$  steps

# SORTING ALGORITHMS

- Given a list, put it into some order
  - Numerical, lexicographic, etc.
- We will see two types
  - Bubble sort
  - Insertion sort

## ALGORITHM 4: BUBBLE SORT

- One of the most simple sorting algorithms
  - Also one of the least efficient
- It takes successive elements and “bubbles” them up the list

**procedure** bubble\_sort ( $a_1, a_2, \dots, a_n$ )

**for**  $i := 1$  **to**  $n-1$

**for**  $j := 1$  **to**  $n-i$

**if**  $a_j > a_{j+1}$

**then** interchange  $a_j$  and  $a_{j+1}$

{  $a_1, \dots, a_n$  are in increasing order }

## BUBBLE SORT RUNNING TIME

```
for  $i := 1$  to  $n-1$   
  for  $j := 1$  to  $n-i$   
    if  $a_j > a_{j+1}$   
      then interchange  $a_j$  and  $a_{j+1}$ 
```

- Outer for loop does  $n-1$  iterations
- Inner for loop does
  - $n-1$  iterations the first time
  - $n-2$  iterations the second time
  - ...
  - 1 iteration the last time
- Total:  $(n-1) + (n-2) + (n-3) + \dots + 2 + 1 = (n^2-n)/2$ 
  - We can say that's “about”  $n^2$  time



# INSERTION SORT RUNNING TIME

```

for  $j := 2$  to  $n$  begin
     $i := 1$ 
    while  $a_j > a_i$ 
         $i := i + 1$ 
     $m := a_j$ 
    for  $k := 0$  to  $j-i-1$ 
         $a_{j-k} := a_{j-k-1}$ 
     $a_i := m$ 
end {  $a_1, a_2, \dots, a_n$  are sorted }
  
```

- Outer for loop runs  $n-1$  times
- In the inner for loop:
  - Worst case is when the while keeps  $i$  at 1, and the for loop runs lots of times
- Total is  $1 + 2 + \dots + n-2 = (n-1)(n-2)/2$ 
  - We can say that's “about”  $n^2$  time

# COMPARISON OF RUNNING TIMES

- Searches
  - Linear:  $n$  steps
  - Binary:  $\log_2 n$  steps
  - Binary search is about as fast as you can get
- Sorts
  - Bubble:  $n^2$  steps
  - Insertion:  $n^2$  steps
  - There are other, more efficient, sorting techniques
    - In principle, the fastest are heap sort, quick sort, and merge sort
    - These each take take  $n * \log_2 n$  steps

